

# Uncertainty Analysis for Complex Systems: Algorithms and Challenges

Dongbin Xiu

Department of Mathematics, Purdue University



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## (Re-)Formulation of PDE: Input Parameterization

$$\frac{\partial u}{\partial t}(t, x) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Goal:** To characterize the random inputs by a set of random variables
  - Finite number
  - Mutual independence
- **If inputs == parameters**
  - Identify the (smallest) independent set
  - Prescribe probability distribution
- **Else if inputs == fields/processes**
  - Approximate the field by a function of finite number of RVs
  - Well-studied for Gaussian processes
  - Under-developed for non-Gaussian processes
  - Examples: Karhunen-Loeve expansion, spectral decomposition, etc.

$$a(x, \omega) \approx \mu_a(x) + \sum_{i=1}^d \tilde{a}_i(x) Z_i(\omega)$$

# The Reformulation

- **Stochastic PDE:**

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \text{ boundary/initial conditions}$$

- **Solution:**  $u(t, x, Z) : [0, T] \times \bar{D} \times \mathbb{R}^{n_z} \rightarrow \mathbb{R}$

- Uncertain inputs are characterized by  $n_z$  random variables  $Z$

- Probability distribution of  $Z$  is prescribed

$$F_Z(s) = \Pr(Z \leq s), \quad s \in \mathbb{R}^{n_z}$$



Non-trivial task

# Generalized Polynomial Chaos (gPC)

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \text{ boundary/initial conditions}$$

• **Focus on dependence on Z:**  $u(\bullet, Z) : \mathbb{R}^{n_z} \rightarrow \mathbb{R}$

•  **$N^{\text{th}}$ -order gPC expansion:**

$$u_N(t, x, Z) \triangleq \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z), \quad \# \text{ of basis} = \binom{n_z + N}{N}$$

• **Orthogonal basis:**  $\int \Phi_i(Z) \Phi_j(Z) \rho(Z) dZ = \delta_{ij}$

• **Basis functions:**

- Hermite polynomials: seminal work by *R. Ghanem*
- General orthogonal polynomials (*Xiu & Karniadakis, 2002*)

• **Properties:**

- Rigorous mathematics
- High accuracy, fast convergence
- Curse-of-dimensionality

• **Numerical Approaches:**

- Galerkin vs. collocation

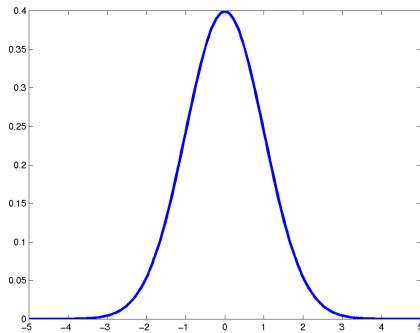
# gPC Basis

▪ **Expectation:**

$$\mathbb{E}(g(Z)) = \int_{\mathbb{R}} g(z)\rho(z) dz$$

▪ **Orthogonality:**

$$\int \Phi_i(z)\Phi_j(z)\rho(z) dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$$

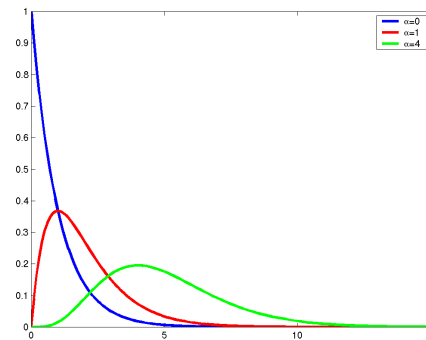


Gaussian distribution

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



Hermite polynomial

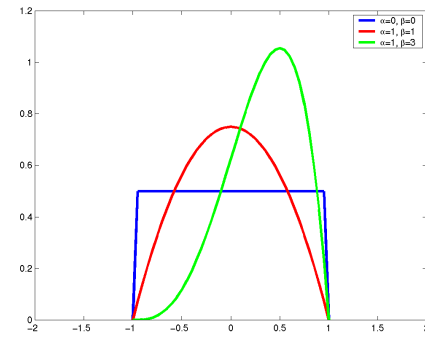


Gamma distribution

$$\int_0^{\infty} \Phi_i(z)\Phi_j(z)e^{-z} dz = \delta_{ij}$$



Laguerre polynomial



Beta distribution

$$\int_{-1}^1 \Phi_i(z)\Phi_j(z) dz = \delta_{ij}$$



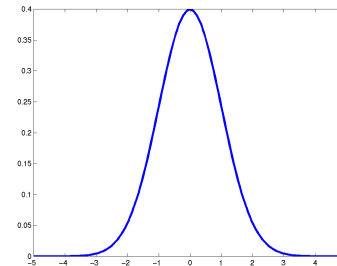
Legendre polynomial

# gPC Basis: the Choices

▪ **Orthogonality:** 
$$\int \Phi_i(z)\Phi_j(z)\rho(z) dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$$

▪ **Example: Hermite polynomial**

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



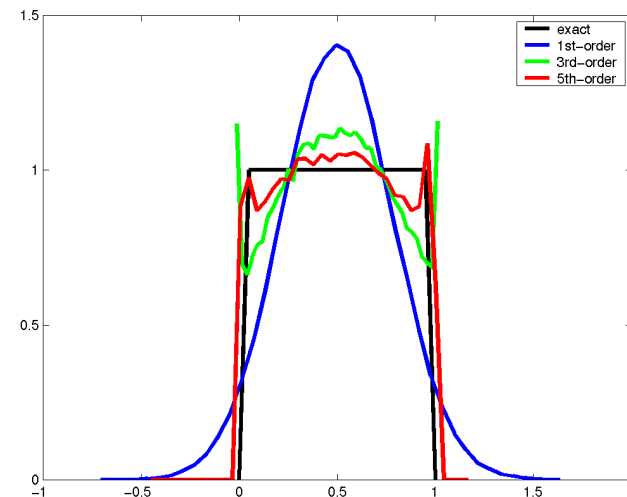
▪ **The polynomials:**  $Z \sim N(0,1)$

$$\Phi_0 = 1, \quad \Phi_1 = Z, \quad \Phi_2 = Z^2 - 1, \quad \Phi_3 = Z^3 - 3Z, \quad \dots$$

▪ **Approximation of arbitrary random variable:** Requires  $L^2$  integrability

▪ **Example:** Uniform random variable

- Convergence
- Non-optimal
- First-order Legendre is exact



# Stochastic Galerkin

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Galerkin method:** Seek

$$u_N(t, x, Z) \triangleq \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z)$$

Such that

$$\mathbb{E} \left[ \frac{\partial u_N}{\partial t}(t, x, Z) \Phi_{\mathbf{m}}(Z) \right] = \mathbb{E} \left[ \mathcal{L}(u_N) \Phi_{\mathbf{m}}(Z) \right], \quad \forall |\mathbf{m}| \leq N$$

- **The result:**
  - Residue is orthogonal to the gPC space
  - A set of deterministic equations for the coefficients
  - The equations are usually coupled – requires new solver

## Stochastic Galerkin: An Example

- Equation : 
$$\frac{du}{dt} = -k(Z)u, \quad u|_{t=0} = u_0. \quad k(Z) = \sum_{i=0}^N k_i \Phi_i(Z)$$

$k(Z)$  is the decaying coefficient with a given probability distribution.

- Seek gPC approximation :

$$v_N(t, Z) = \sum_{i=0}^N \hat{v}_i(t) \Phi_i(Z)$$

- Galerkin equation :

$$\frac{d\hat{v}_k}{dt} = -\sum_{i=0}^N \sum_{j=0}^N e_{ijk} k_i \hat{v}_j, \quad k = 0, 1, 2, \dots, N$$

$$e_{ijk} = \int \Phi_i(z) \Phi_j(z) \Phi_k(z) \rho(z) dz$$

- Computational complexity:  $(N+1)$  coupled deterministic ODEs

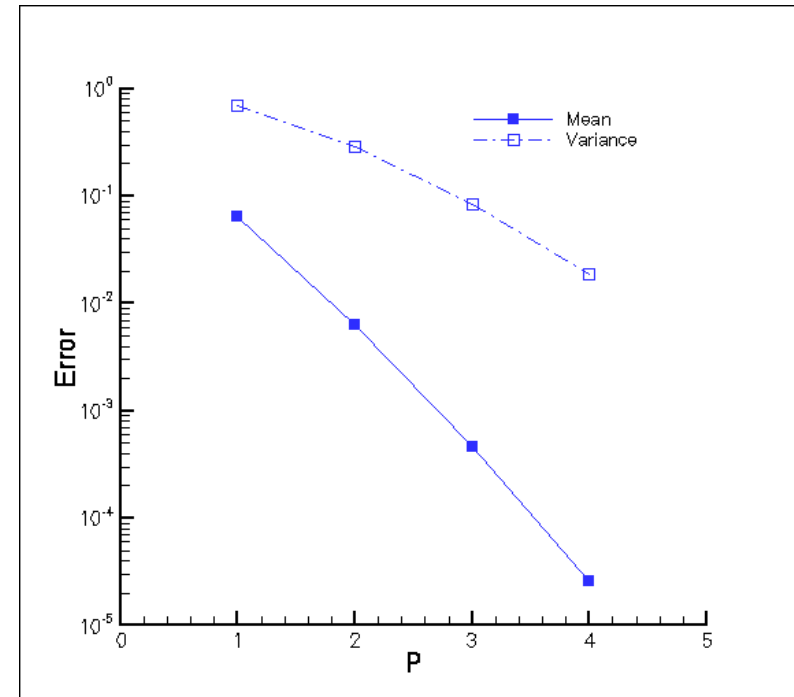


# Computational Efficiency

- $du/dt = -k u$ ,  $u(t=0)=1$
- $k$  is a **Gaussian** random variable :

$$\text{PDF: } f_k(x) = \frac{1}{\sqrt{2p}} e^{-\frac{x^2}{2}}$$

- **4<sup>th</sup>-order Hermite expansion**



Error	Monte Carlo Method (# of realizations)	Generalized Polynomial Chaos (# of expansion terms)	Speed-up factor
4%	100	1	100
1.1%	1,000	2	500
0.05%	9,800	3	3,267

# Stochastic Collocation

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \text{ boundary/initial conditions}$$

- **Collocation:** To satisfy governing equations at selected nodes
    - Allow one to use existing deterministic codes repetitively
- 

- **Sampling:** (solution statistics only)
    - Random (Monte Carlo)
    - Deterministic (lattice rule, tensor grid, cubature)
- 

- **Stochastic collocation:** To construct **polynomial approximations**
  - Node selection is critical to efficiency and accuracy
  - More than sampling

**Definition:** Given a set of nodes and solution ensemble, find  $p(Z)$  in a proper polynomial space, such that  $p \approx u$  in a proper sense.

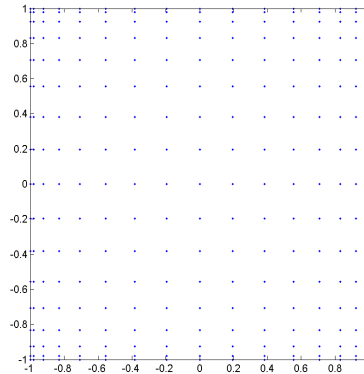
# Stochastic Collocation: Interpolation

- **Lagrange interpolation:**

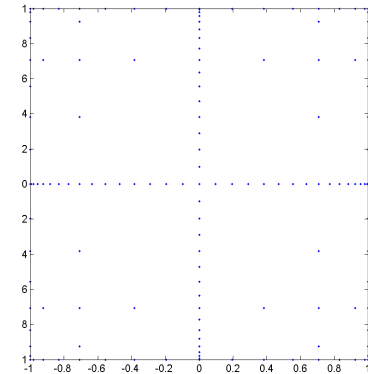
- Let  $z_j$  be the nodes and  $u(z_j)$  be solution, then Lagrange interpolation

$$p(z) = \sum_{j=1}^Q u(z_j) L_j(z) \quad L_i(z_j) = \delta_{ij}, \quad 1 \leq i, j \leq N_p$$

- Difficult for unstructured grids.
- Dimension-by-dimension space filling



Tensor grids: inefficient



Sparse grids: more efficient

- **Matrix inversion:**

$$p(Z) = \sum_{i=1}^M c_i \Phi_i(Z)$$

$$p(z_j) = \sum_{i=1}^M c_i \Phi_i(z_j) = f_j \quad \Rightarrow \quad \mathbf{A} \mathbf{c} = \mathbf{f}$$

**Vandermonde matrix:**  $\mathbf{A} = (a_{jk}) = (\Phi_k(z_j)), \quad j = 1, \dots, N_p, \quad k = 1, \dots, M$

# Stochastic Collocation: Non-interpolating

- **Regression type:**

$$\min \|\mathbf{A}\mathbf{c} - \mathbf{f}\|$$

- Over-determined system: least-square type
- Under-determined system:  $l_1$ -minimization, compressive sampling, etc.

- **Discrete projection:**

$$\mathbb{P}_N u = \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z)$$

$$\hat{u}_{\mathbf{k}} = \mathbb{E}[u(Z)\Phi_{\mathbf{k}}(Z)] = \int u(z)\Phi_{\mathbf{k}}(z)\rho(z) dz$$

$$\approx \sum_{j=1}^{N_p} u(z_j)\Phi_{\mathbf{k}}(z_j)w_j$$

# Stochastic Computation: The Landscape

- **Realistic Large-scale Complex Systems:**

- Complex physics → highly nonlinear systems
- Large number of random variables
- (Extremely) time consuming simulations
- Legacy codes (nearly impossible to re-write)

- **Stochastic Galerkin:**

- Difficult to implement
- Good mathematical properties

- **Stochastic collocation is more proper:**

- Easy to implement → virtually no coding effort
- Nonlinearity poses no additional difficulties

- **Easy implementation:**

1. Choose a set of nodes,  $Z_j, j=1, \dots, N_p$ .
2. Run deterministic simulation at each node  $Z_j$ .
3. Construct polynomial approximation (surrogate/response surface).

# Stochastic Computation: Challenges

- **Curse-of-Dimensionality:**

- Number of simulations grows (too) fast with dimensionality
- Current approaches:
  - Adaptive (sparse) grid
  - “Sparser” grids
- Significantly “delayed” but far from satisfactory
  - A rather extreme (but not uncommon) scenario:  
*“What if I have 30 random inputs but can only afford 10 simulations?”*

- **Do we know all the probability distributions?**

- In many practical systems, we do not → Epistemic uncertainty
- Very few studies
- (Probably) the first numerical approach: *Jakeman, et al, JCP 2010*

- **Multi-physics, multi-scale systems**

# Stochastic Computation: “Useful” Algorithms

- “Useful” UQ algorithms need to target ....
  - *Realistic Large-scale Complex Systems:*
    - *Complex physics → highly nonlinear systems*
    - *Large number of random variables*
    - *(Extremely) time consuming simulations*
    - *Legacy codes (nearly impossible to re-write)*
- **More development of “capability-based” UQ**
  - To make UQ algorithms with certain capability/accuracy more efficient
  - For example: adaptive refinement
- **In need of “capacity-based” UQ**
  - To design the “best” method for a given simulation capacity
  - For example:
    - “*What if I have 30 random inputs but can only afford 10 simulations?*”
    - Rephrase: “*Assume we can afford 10 simulations, what can we achieve?*”

# Summary

- **Uncertainty Analysis:** To provide improved prediction
  - Input characterization
  - Uncertainty propagation
  - Post processing
- **Generalized polynomial chaos (gPC)**
  - Multivariate approximation theory
- **Active directions:**
  - Compressive sampling
  - Adaptive algorithms
  - Model-form uncertainty
  - Utilization of data: data assimilation, inference, etc.
  - etc, etc, etc...
- **What about visualization?**
  - **Lack of dialogue between the UQ and Viz communities**